

More on Java Synchronization

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Outline

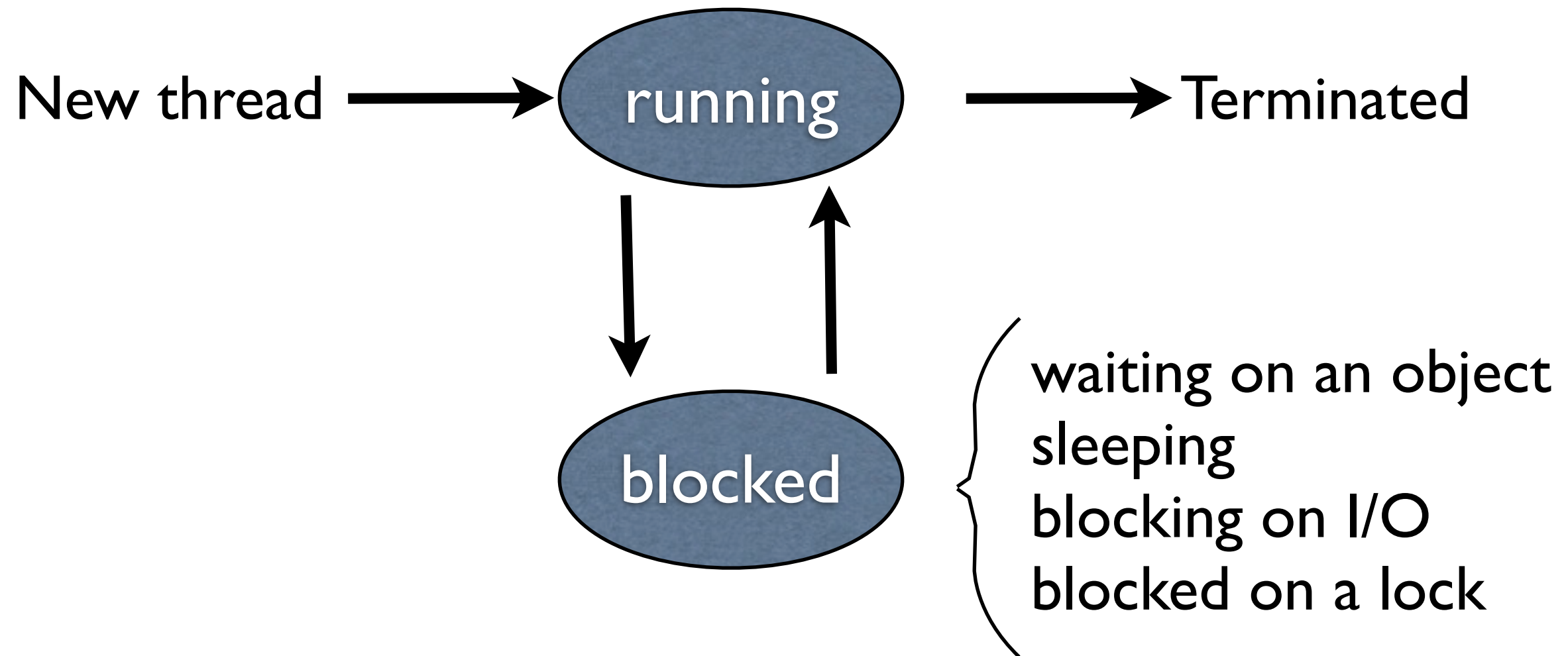
- Basics of Thread Operations (new stuff)
- Thread States
- Thread Scheduling
- Threads and I/O
- Producer/Consumer Pattern
- wait()/notify()
- Thread-Safe Data Structures
- Processing Lab

The Basics

- **Threads**

- `run()/start()`
- `yield()`
- `sleep()`
- `join()`
- `wait()` and `notify()`, and also `notifyAll()`

States of a Thread



How to get the state?

<http://docs.oracle.com/javase/1.5.0/docs/api/java/lang/Thread.State.html>

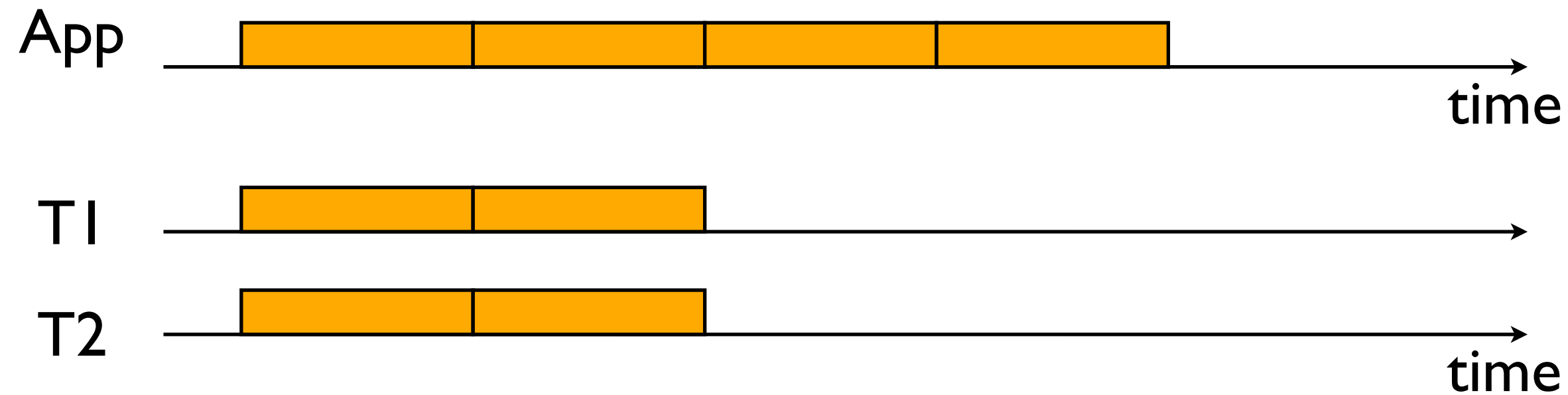
- NEW
- RUNNABLE
- BLOCKED
- WAITING
- TIME_WAITING
- TERMINATED

`getState()`

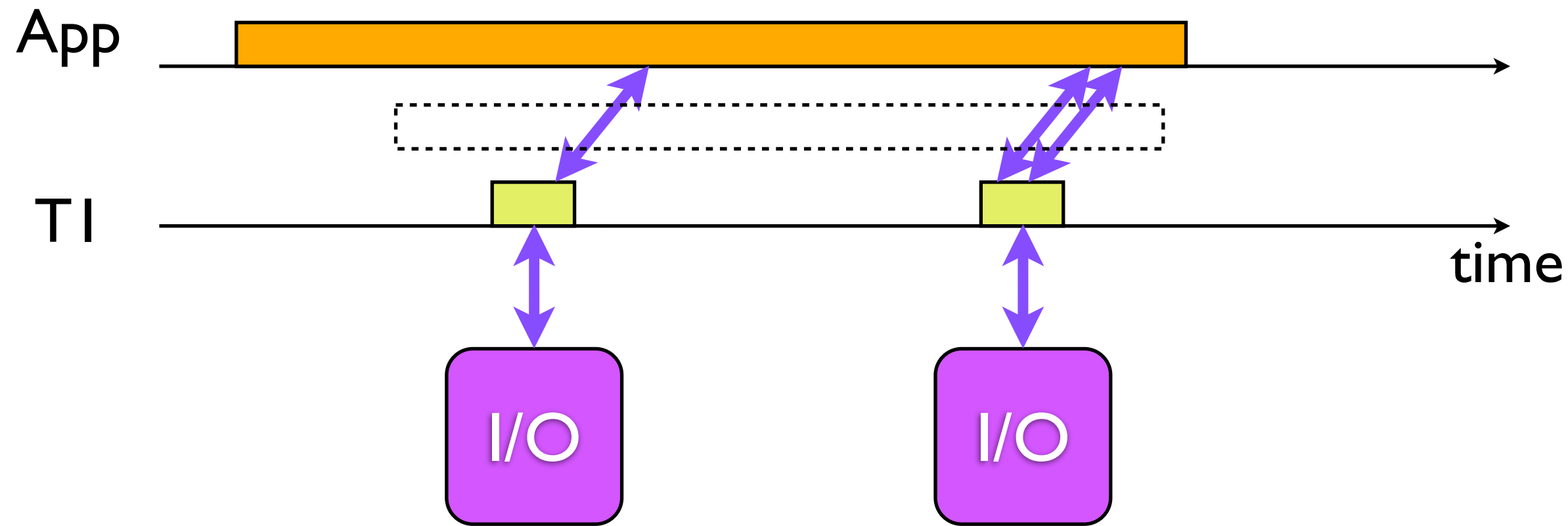
Thread Scheduling

- What is the policy?
 - Java doc says: Implemented in the JVM, preemptive, based on priority. (No mention of time-slices.)
 - 1 = low priority, 5 = main, 10 = high priority
 - `getPriority()` & `setPriority()`
 - However, most OS implement **time-slices** (quanta), roughly 1ms, **preemptive**, and **round-robin** ==> JVMs do the same

Threads good not only for speedup



Threads good not only for speedup But also to simplify code



Important Concepts

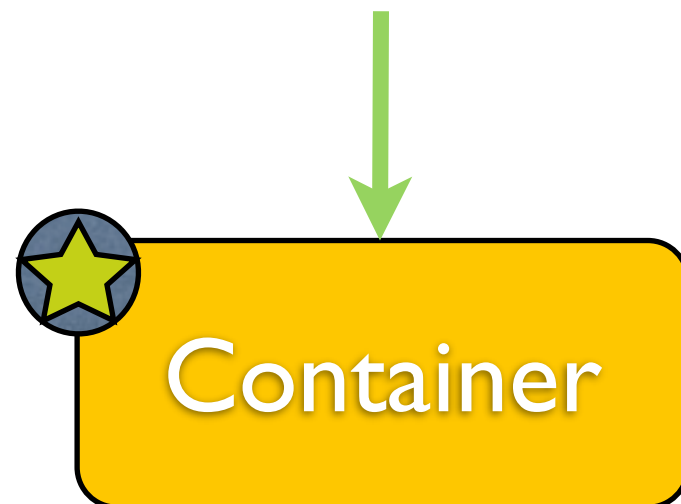
- CPU Bound Processes/Threads
- I/O Bound Processes/Threads

Time Scale

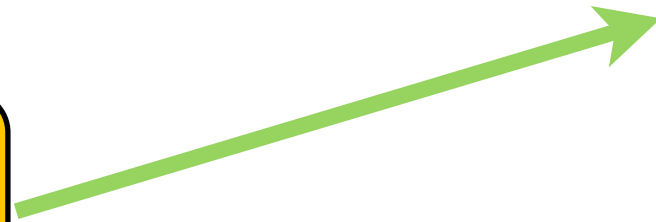
- Why I/O recognizing I/O-bound process is important
 - CPU cycle: 1 ns
 - RAM cycle: 100-500 ns
 - Disk access = seek + latency
 - seek = 1 ms
 - latency = 1/2 rotation, at 7,000 RPM
- Question: How long does the processor wait for data from disk?

Wait/Notify Example (Producer-Consumer)

```
// producer code  
synchronized( lock ) {  
    while ( !container.isEmpty() ) {  
        try {  
            lock.wait();  
        } catch (InterruptedException e)  
        {}  
    }  
    container.put( newItem );  
}
```



```
//consumer code  
synchronized( lock ) {  
    if ( !container.isEmpty() ) {  
        item = container.getItem();  
        consume( item );  
        lock.notify();  
    }  
}
```



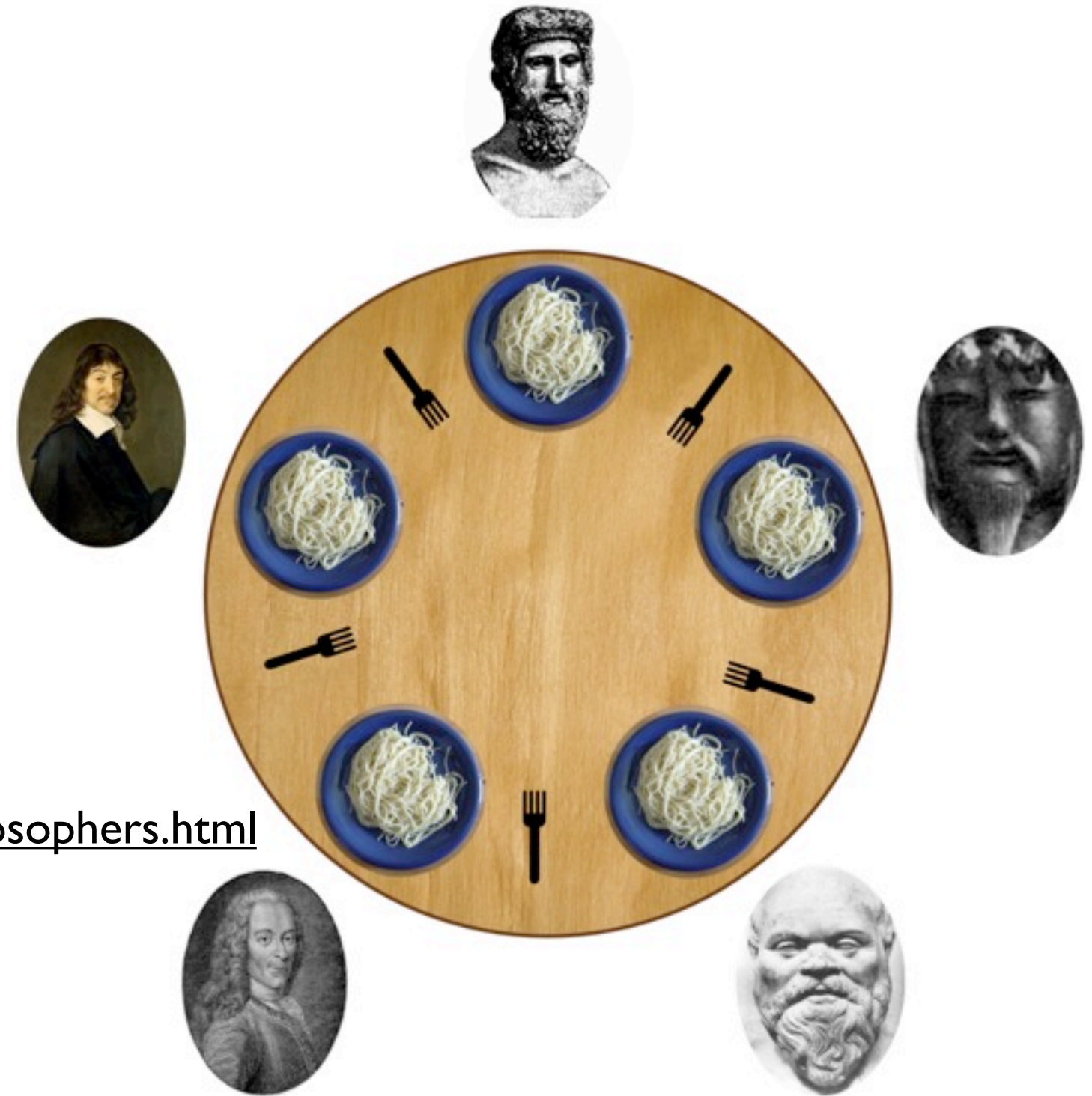
Beware of Deadlocks!

- The Dining-Philosophers Problem

http://en.wikipedia.org/wiki/Dining_philosophers_problem

- The Applet

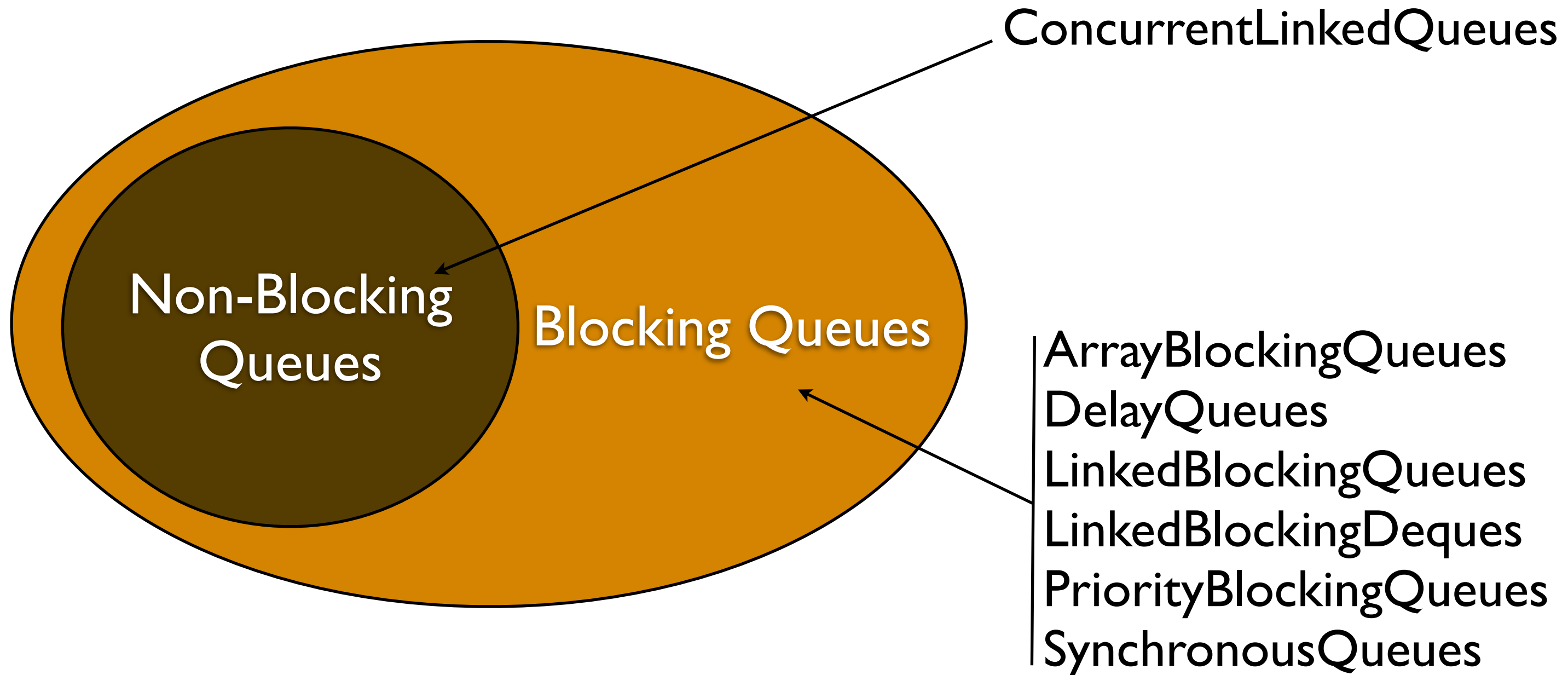
<http://elvis.rowan.edu/~hartley/ConcProgJava/Applets/diningPhilosophers.html>



Rule #1 for Preventing Deadlocks

- Grab all the data-structures that you need first
- If you can't, release them all
- Wait a random amount of time and try again

Thread-Safe Data Structures



ConcurrentLinkedQueue

Method Summary	
boolean	<code>add(E e)</code> Inserts the specified element at the tail of this queue.
boolean	<code>contains(Object o)</code> Returns <code>true</code> if this queue contains the specified element.
boolean	<code>isEmpty()</code> Returns <code>true</code> if this queue contains no elements.
<code>Iterator<E></code>	<code>iterator()</code> Returns an iterator over the elements in this queue in proper sequence.
boolean	<code>offer(E e)</code> Inserts the specified element at the tail of this queue.
<code>E</code>	<code>peek()</code> Retrieves, but does not remove, the head of this queue, or returns <code>null</code> if this queue is empty.
<code>E</code>	<code>poll()</code> Retrieves and removes the head of this queue, or returns <code>null</code> if this queue is empty.
boolean	<code>remove(Object o)</code> Removes a single instance of the specified element from this queue, if it is present.
int	<code>size()</code> Returns the number of elements in this queue.
<code>Object[]</code>	<code>toArray()</code> Returns an array containing all of the elements in this queue, in proper sequence.

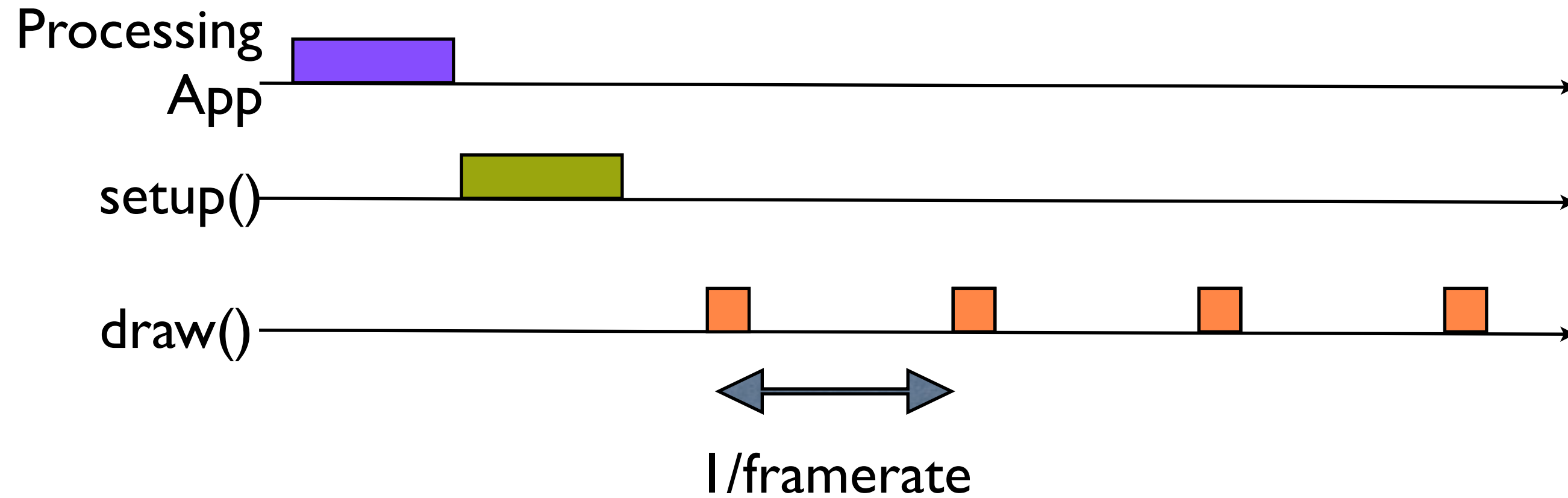
ArrayBlockingQueue

Method Summary	
void	clear() Removes all of the elements from this collection.
boolean	contains(Object o) Returns true if this collection contains the specified element.
int	drainTo(Collection<? super E> c) Removes all available elements from this queue and adds them into the given collection.
int	drainTo(Collection<? super E> c, int maxElements) Removes at most the given number of available elements from this queue and adds them into the given collection.
Iterator<E>	iterator() Returns an iterator over the elements in this queue in proper sequence.
boolean	offer(E o) Inserts the specified element at the tail of this queue if possible, returning immediately if this queue is full.
boolean	offer(E o, long timeout, TimeUnit unit) Inserts the specified element at the tail of this queue, waiting if necessary up to the specified wait time for space to become available.
E	peek() Retrieves, but does not remove, the head of this queue, returning null if this queue is empty.
E	poll() Retrieves and removes the head of this queue, or null if this queue is empty.
E	poll(long timeout, TimeUnit unit) Retrieves and removes the head of this queue, waiting if necessary up to the specified wait time if no elements are present on this queue.
void	put(E o) Adds the specified element to the tail of this queue, waiting if necessary for space to become available.
int	remainingCapacity() Returns the number of elements that this queue can ideally (in the absence of memory or resource constraints) accept without blocking.

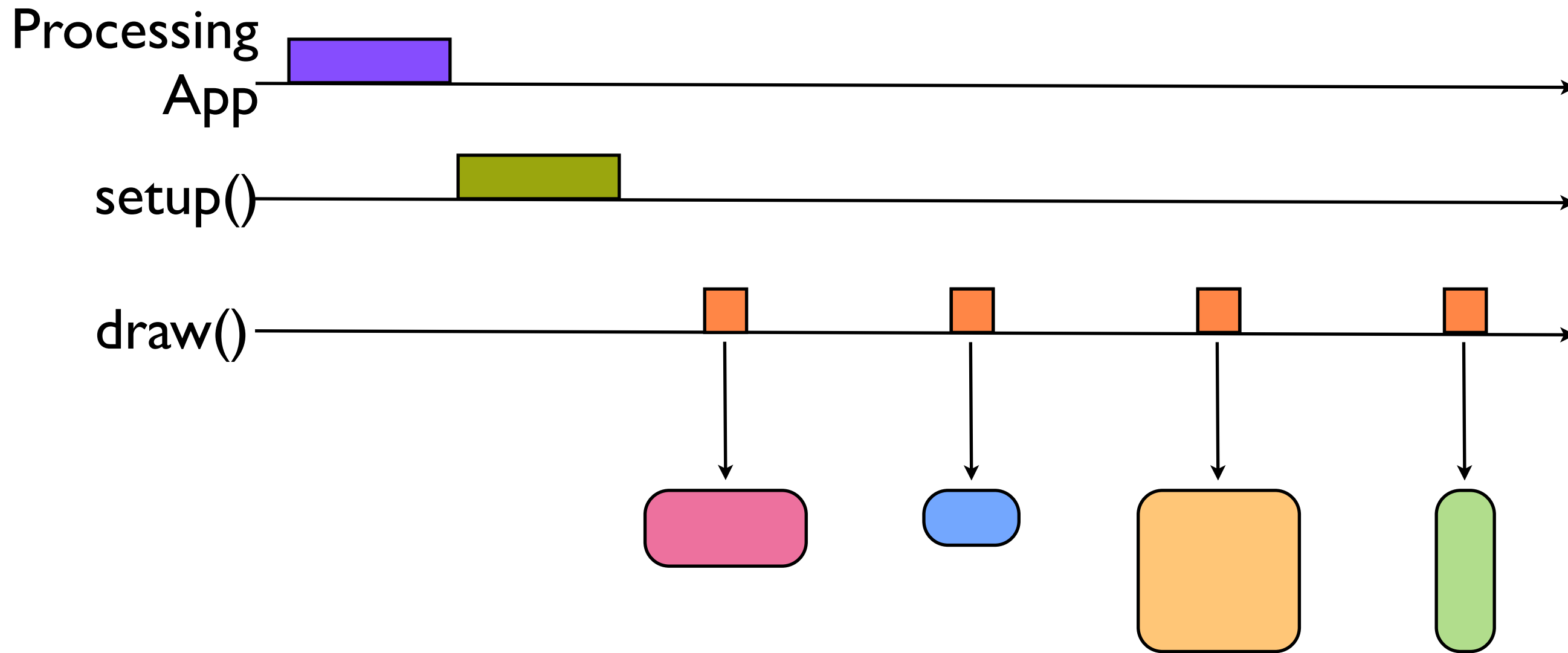
ArrayBlockingQueue (cont'd)

int	remainingCapacity() Returns the number of elements that this queue can ideally (in the absence of memory or resource constraints) accept without blocking.
boolean	remove(Object o) Removes a single instance of the specified element from this collection, if it is present (optional operation).
int	size() Returns the number of elements in this queue.
E	take() Retrieves and removes the head of this queue, waiting if no elements are present on this queue.
Object[]	toArray() Returns an array containing all of the elements in this collection.
<T> T[]	toArray(T[] a) Returns an array containing all of the elements in this collection; the runtime type of the returned array is that of the specified array.
String	toString() Returns a string representation of this collection.

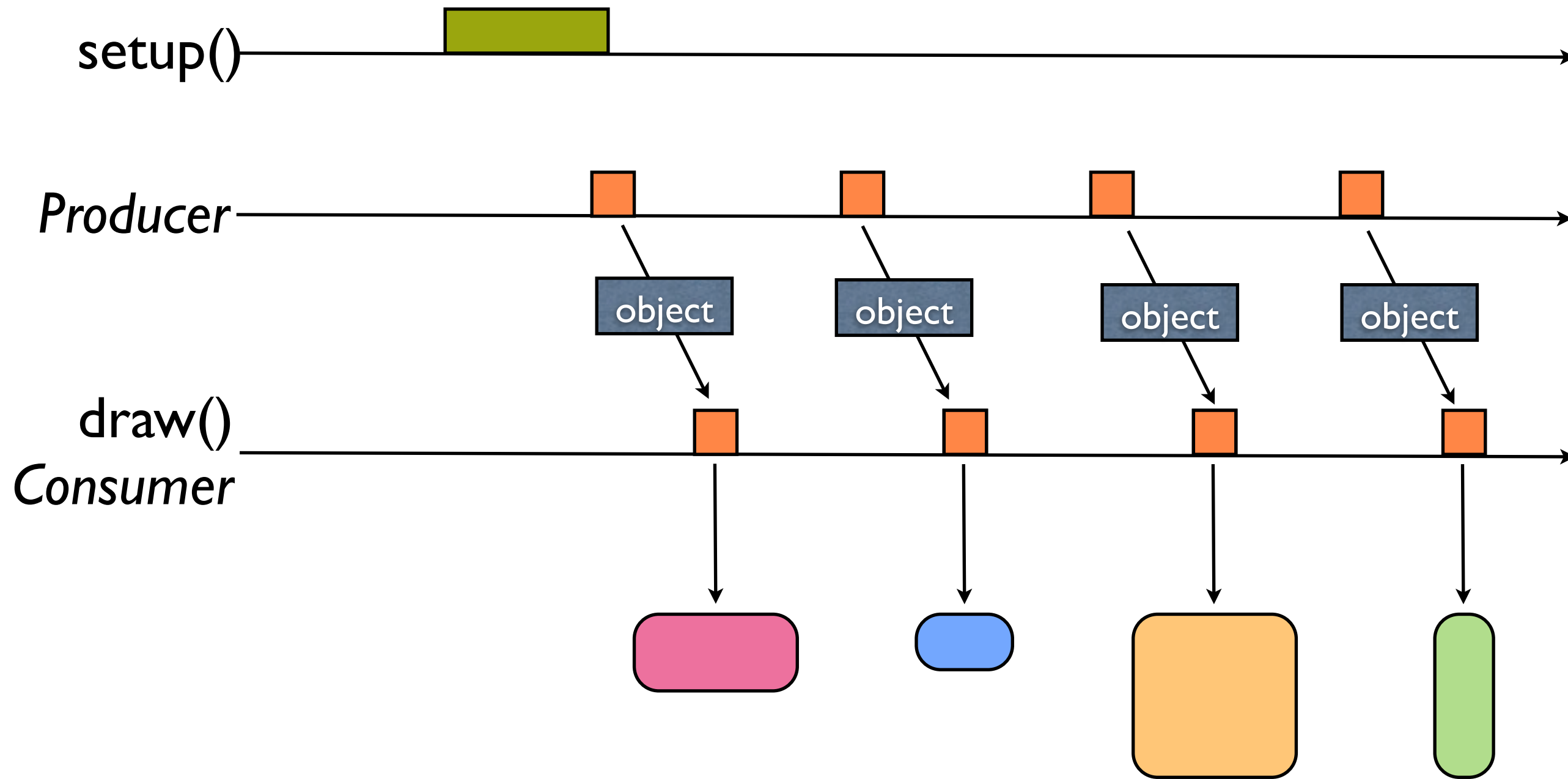
Processing Lab



Processing Lab: Version I



Processing Lab: Version I



Processing Lab: Version 2

